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## **Modelling Time-Constrained Software Development**

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**This paper is circulated for discussion purposes only and its contents should be  
considered preliminary.**

## **Abstract**

Commercial pressures on time-to-market often require the development of software in situations where deadlines are very tight and non-negotiable. This type of development can be termed ‘time-constrained software development.’

The need to compress development timescales influences both the software process and the way it is managed. Conventional approaches to modelling tend to treat the development process as being linear, sequential and static. Whereas, the processes used to achieve timescale compression in industry are iterative, concurrent and dynamic. That is, they replace the notion of ‘right-first-time’ with one of ‘right-on-time.’

In this paper we propose a new modelling technique, called Capacity-Based Scheduling (CBS), to control risk across a portfolio of time-constrained projects. We show how schedule constraints can be modelled in order to predict the consequences of alternative plans and control schedule risk across a portfolio of time-constrained projects.

## 1. Introduction

The pressures on software development timescales are growing. Improvements, such as the use of commercial off-the-shelf technology, have caused both the costs and timescales associated with hardware to fall, thus making software a critical path component in the provision of many products and services (Sims 1997).

Organisations must therefore continuously seek ways to improve their lead-time capability and achieve more with diminishing resources. These improvements must come from using: the right people, processes and tools; eliminating unnecessary work; being right-first-time; or by doing many things at once (Parkinson 1996). The ‘*time-to-market*’ principle assumes that savings or benefits outside of the process can outweigh any additional costs of compressed timescales.

The demands on lead-time compression are increasingly being met by the application of *evolutionary development* lifecycles. These overcome the problems of conventional lifecycles, typified by the Waterfall model (Royce 1970), that imply a sequential once-through approach to development. These models assume that software can be developed ‘right-first-time’ and that lead-times are sufficiently long to proceed in a stepwise manner. Instead, evolutionary lifecycles such as the Spiral model (Boehm 1988) recognise the need for software artefacts to evolve over time in a controlled risk-driven manner. These lifecycles have effectively replaced the notion of ‘*right-first-time*’ with ‘*right-on-time*.’

In the remainder of this paper, we explore the limitations of current predictive models and propose a new approach to modelling time-constrained development. We start by identifying the problems of concurrency and iteration (Section 2) and conventional models (Section 3). We then propose a new approach called Capacity-Based Scheduling (CBS) to help managers to plan and control risk across a portfolio of projects (Section 4). Finally, we describe an implementation of CBS using a primitive model to explain its operations by way of example (Section 5).

## 2. Concurrency and Iteration

The two essential elements of evolutionary lifecycles are iteration and concurrency. *Iteration* is the repetition of development activities to deliver increments of product functionality at pre-planned intervals. Incremental development, or staged-delivery, balances progress and early feedback against the overheads of repeating the process a number of times. *Concurrency* is the simultaneous performance of development activities between projects, product deliveries, development phases and individual tasks.

The extent of concurrency and iteration therefore distinguishes time-constrained software development from more conventional processes. These differences can be illustrated as three 'modes' of software development as illustrated in Figure 1.

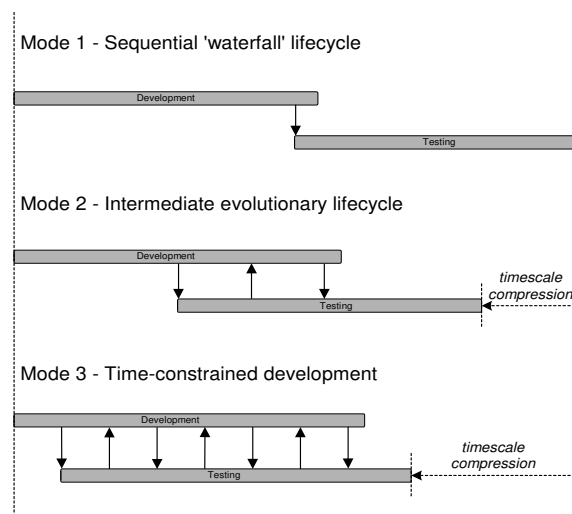


Figure 1: Three Modes of Software Development

*Mode 1* represents a sequential 'waterfall' model (Royce 1970) of software development that lays development activities end-to-end. The intrinsic assumption that software can be developed right-first-time is, however, unrealistic and the resulting timescale is commercially impractical for many organisations (Boehm 1988).

*Mode 2* represents an intermediate evolutionary lifecycle using a limited amount of concurrency and iteration. This allows feedback in the software process to detect and fix

problems whilst, in theory, reducing the overall development lead-time. This approach is typical of conventional software development (Parnas and Clements 1986).

*Mode 3* represents the extreme form of evolutionary development that we have observed in industry. This exploits high levels of concurrency and staged-delivery to meet the time constraints imposed on the software process, and is based on three assumptions. First, that concurrency reduces overall development lead-times. Second, that iteration gives rise to improvements in product maturity. Third, that the benefits of reduced lead-times outweigh the costs of instability and rework.

The existence and effects of lifecycle concurrency and iteration are, however, largely implicit in the literature and are treated as an issue of project management rather than an integral part of the software lifecycle (Blackburn and Hoedemaker 1996). Conventional approaches to measurement, modelling and management treat the development process as being linear, sequential and static. Whereas, the industrial practices used to achieve timescale compression in industry are iterative, concurrent and dynamic.

An important observation is that a schedule delay, or a change in the quantity of work performed at any point in these lifecycles, has the potential to cause increased allocation of resource that can affect the performance of later phases. This dynamic behaviour has significant implications for the way in which software development is modelled (Abdel-Hamid and Madnick 1991; Lehman 1995). It follows that we need to reassess existing predictive models for their suitability in time-constrained environments.

### **3. Conventional Predictive Models**

Conventional predictive models take the form shown in Figure 2. For example, COCOMO takes the user's predictions of system size and, using 15 productivity drivers (e.g. programming languages, programmer experience), estimates the level of effort required in person-months (Boehm 1981). A planning tool, such as Microsoft Project, is then used to arrange schedules and level resource demands in the context of other parallel projects.

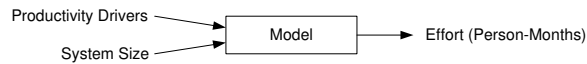


Figure 2: Conventional Predictive Models of Software Development

In practice, the predictive accuracy of these models is poor even for simple development projects. As Kitchenham observes “*There is no evidence that estimation models can do much better than get within 100% of the actual effort during requirements specification and 30% of the actual effort prior to coding*” (Kitchenham 1998). The models are even less appropriate for managing time-constrained software development for a number of reasons.

First, in time-constrained development (i) the timescales of the project, and its major increments, are largely fixed and non-negotiable, and (ii) the total level of resource remains, on the whole, quite stable and predictable. Since time and resource (and thus available effort) are ‘known’ variables, having them as outputs of a predictive model is therefore of little benefit to planners.

Second, as time-constrained development is a dynamic feedback system, estimating and planning cannot be treated as separate activities. The ‘point-estimates’ given by conventional estimation models therefore give little indication of the risks involved in accepting these constraints because process behaviour is dynamically dependent on the way the process is planned and controlled, i.e. “*a different estimate creates a different project*” (Abdel-Hamid 1989).

The intrinsic problems of prediction mean that we need to rethink its rôle in the management of software development. As Kitchenham observes: “*senior managers and project managers need to concentrate more on managing estimate risk than looking for a magic solution to the estimation problem*” (Kitchenham 1998).



## 4. A Model of Time-Constrained Software Development

### 4.1 Overview of Capacity Based Scheduling

We have developed a new model of time-constrained development, called *Capacity Based Scheduling (CBS)*, which addresses the problems of the lead-time approach. The basic concept of Capacity-Based Scheduling is to reverse the inputs and outputs of current predictive models (

Figure 3).

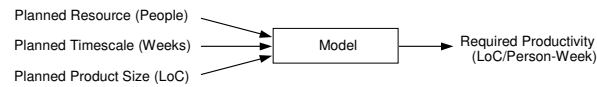


Figure 3: Capacity-Based Scheduling – Concept

By making the constraints on resources and timescales explicit inputs, we can reason about the nature of the capabilities required for their achievement. Critically, we are modelling the constraints on the process in order to compare the *Required Productivity* against the *Actual Productivity* of past projects, and thus indicate the relative ‘risk’ of alternative plans.

The principles behind the model are best explained by example. Consider two planned deliveries of code,  $A_1$  and  $A_2$ . Delivery  $A_1$  needs 1000 Lines of Code (LoC) to be written and delivery  $A_2$  needs 500 LoC. The work can start on each delivery immediately and the delivery deadline for each is 10 weeks from now. Supposing we have a staff of six programmers, how should we allocate staff between the parallel deliveries?

At time 0, delivery  $A_1$  requires 100 LoC to be produced each week until the deadline and  $A_2$  requires 50 LoC to be produced each week until the deadline. An obvious and well-motivated approach to allocating staff to the tasks would be to allocate four programmers to  $A_1$  and two programmers to  $A_2$ . This balances the load between the parallel activities; no programmer is under-stressed whilst another is overstressed. In other words, this is *proportional allocation according to need*.

Thus, after the first week with the proportional allocation above, the four programmers working on  $A_1$  deliver (together) 100 LoC and the two on  $A_2$  deliver 50 LoC. If there are no other deliveries starting after the first week then the required rates of progress for  $A_1$  and  $A_2$  are now  $900/9=100$  LoC per week and  $450/9=50$  LoC per week, i.e. the same as before. The same approach can likewise be used for any number of concurrent deliveries; if, for example, another delivery,  $A_3$ , started after the first week then we would need to take it into account in deciding proportional allocations for Week 2.

The *Required Productivity* at any stage can be compared with historical performance and judgements made. For example, requiring programmers to have a very high productivity now could affect the amount of work to be carried out in a later delivery since stressed programmers might introduce more defects. By modelling planning constraints and assumptions, the CBS approach highlights what meeting the deadlines actually means for the staff carrying out the tasks.

## 4.2 A Formal Description of the Primitive Model

The basic structure of our primitive model is presented in Figure 4 that shows the main variables and their causal relationships as arrows. The variables are calculated by dynamic (time-based) equations that describe the relationship between inputs and outputs.

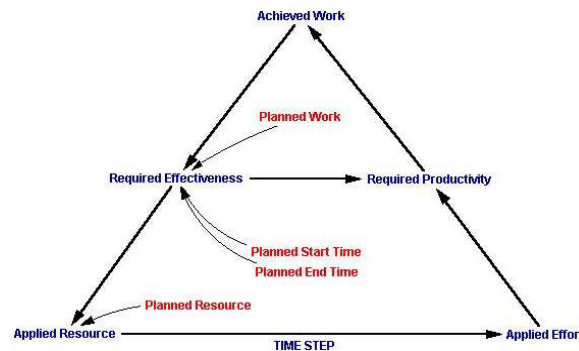


Figure 4: Capacity-Based Scheduling – Primitive Model

The inputs to the model are the *Planned Work* (LoC) to be produced, the *Planned Resource* (people) in each phase team, the *Planned Start Time* (time), and the *Planned End Time* (time) of each phase delivery. The outputs are the calculated values of *Required Effectiveness*

(LoC/week), the *Applied Resource* (people), the *Applied Effort* (person-weeks), the *Required Productivity* (LoC/person-week), and the *Achieved Work* (LoC), over time.

The following steps formally describe one cycle of the model's operation (corresponding to one unit of time). The full model is implemented using the VenSim simulation environment. Here we concentrate on the operations of the primitive model introduced to explain the modelling approach.

### Step 1 – Input the Planned Schedule

There are a number of planned deliveries (e.g.  $D_1, D_2$ ).

- $Deliveries = \{D_1, ..., D_n\}$

Each delivery comprises several phases of the software lifecycle (e.g. *Code, Unit Test*).

- $Phases = \{P_1, ..., P_n\}$

The planning window comprises a set of time points of from 0 to  $T$  with the current time being  $t$ .

- $Time = 0 ... T$
- $time = t$

Each delivery involves planned work required for its completion (e.g. 1000 LoC for  $D_1$  and 500 LoC for  $D_2$ ).

- $Planned\ Work : Phase \times Delivery \rightarrow R$   
(where  $R$  is the set of non-negative numbers)
- Each delivery phase has a planned start and end time.
- $Planned\ Start\ Time : Phase \times Delivery \rightarrow Time$
- $Planned\ End\ Time : Phase \times Delivery \rightarrow Time$

Finally, there is a fixed resource capability available for each phase:

- *Planned Resource Phase: Phase*  $\rightarrow R$   
(i.e. a natural number of people)

## Step 2 – Calculate the Required Effectiveness

After  $t$  time steps a certain amount of work has been performed (*Work Achieved*) for a specific *Phase* ( $p$ ) of each *Delivery* ( $d$ ).

- *Achieved Work : Phase x Delivery x Time*  $\rightarrow R$

Each product is associated with the remaining required effectiveness at time  $t$ .

- *Required Effectiveness : Work x Time*  $\rightarrow R$  (LoC/time) given by
- *Required Effectiveness* ( $p,d,t$ ) = (*Planned Work* ( $p,d,t$ ) – *Achieved Work* ( $p,d,t$ )) / (*Planned End Time* ( $p,d$ )- $t$ )
- if *Planned End Time* ( $p,d$ )  $> t$  and 0 otherwise.

The *Required Effectiveness* shows the total stress on the development team.

## Step 3 – Calculate the Applied Resource

At each time interval ( $t$ ), the *Planned Resource* for each phase ( $p$ ) is allocated among parallel deliveries ( $d$ ) relative to the *Required Effectiveness*.

- *Applied Resource: Phase x Delivery x Time*  $\rightarrow R$
- *Applied Resource* ( $p,d,t$ ) = *Planned Resource*( $p$ )  $\times$  (*Required Effectiveness*( $p,d,t$ )  
 $/ \sum_{j=D_1}^{j=D_n} \text{Required Effectiveness}(p,j,t)$ )

## Step 4 – Calculate the Applied Effort

The *AppliedEffort* is the *AppliedResource* over a time interval ( $t$ ).

- *Applied Effort* :  $Phase \times Delivery \times Time \rightarrow R$
- $Applied\ Effort(p,d,t) = Applied\ Resource(p,d,t) \times t$

### Step 5 – Calculate the Required Productivity

We are interested in how much the *Required Productivity* might have to deviate from known performance in order to meet the deadline.

- *Required Productivity* :  $Phase \times Delivery \times Time \rightarrow R$  given by
- $Required\ Productivity(p,d,t) = Required\ Effectiveness(p,d,t) / Applied\ Effort(p,d,t)$

if  $Applied\ Effort(p,d,t) > 0$  and 0 otherwise.

### Step 6 – Calculate the Achieved Work

The *Achieved Work* in each time-step is assumed to be the amount required.

- *Achieved Work* :  $Phase \times Delivery \times Time \rightarrow R$  given by
- $Achieved\ Work(p,d,t) = Required\ Productivity(p,d,t) \times Applied\ Resource(p,d,t)$

Finally, the *Achieved Work* to-date for a delivery is given by summing all the achievement for that product to date.

- $Achieved\ Work\ To-Date(p,d,t) = \sum_{j=1}^{j=t} Achieved\ Work(p,d,j)$

### Step 7 – Iterate

The model iterates between Steps 2 and 6 until time equals the end time ( $T$ ).

## 5. Model Worked Examples

In this section, three examples are shown to demonstrate the operation of the basic model. The first example, Section 5.1, revisits our example from Section 4.2 to model two concurrent deliveries for a single lifecycle phase (*Code*). The second example, Section 5.2, considers the impact of staging our two concurrent deliveries to start at different points in time. The final example, Section 5.3, uses the model to study the effect of having staged concurrent deliveries across two lifecycle phases (*Code* and *Low Level Test*). In all cases, the goal is to evaluate whether the planned schedule is feasible given external constraints and internal planning assumptions.

### 5.1 Example 1: Two Deliveries, One Phase

The input to the model is the planned delivery schedule consisting of the: *Planned Resource*, *Planned Work*, *Planned Start Time* and *Planned End Time*, for each delivery.

The *Planned Resource* is the number of available resource (people) for each phase. We assume that each phase (e.g. *Code*) contains a fixed pool of staffing resources over the time-period under consideration. In this example, the *Planned Resource* is 6 people.

The *Planned Work* is the total amount of work (LoC) to be performed. We assume that this plan includes both planned new work and an allowance for planned rework. In this example, the values of *Planned Work* are 1,000 LoC and 500 LoC respectively.

The *Planned Start Time* and *Planned End Time* are the planned start and end times (week) of each phase delivery. In this example, both *A1 Code* and *A2 Code* start at Week 0 and end at Week 10 (Figure 5a, situated at the end of this paper).

The model then iteratively calculates the values of *Required Effectiveness*, *Applied Resource*, *Applied Effort*, *Required Productivity* and *Achieved work*.

The *Required Effectiveness* calculates the average work rate (LoC/week) required to meet the deadline. We assume that the coders are 100% effective, i.e. our estimates of size are correct and there is no rework. In this example, the *Required Effectiveness* of delivery *A1* is

therefore 100 LoC per week and delivery A2 is 50 LoC per week for the duration under study (Figure 5b).

The *Applied Resource* is allocated according to the *Required Effectiveness* over time of the two parallel deliveries. Our strategy for allocating resources does not take into account other deliveries with start times in the future but only with the active tasks in the next time step (week). We assume that deliveries share a common resource team that is allocated by managers according to the relative size of the task. In this case, the ratio of *Required Effectiveness* is 100:50, or 2:1, so the *Planned Resource* of 6 people is allocated in the ratio of 4 coders on A1 to 2 coders on A2 (Figure 5c).

The *Applied Effort* is the product of *Applied Resource* and *Time*. We assume a standard working week with no overtime such that total effort is fixed, therefore the *Applied Effort* mirrors the level of *Applied Resource*. In this example, the *Applied Effort* is consistent at 4 person-weeks for A1 and 2 person-weeks for A2 (Figure 5d).

The total *Required Productivity* is equivalent to the *Required Effectiveness* since we assume that progress is actually made evenly at the required rate (i.e. all deadlines are met). In this example, if 25 LoC/person-week is the average *Required Productivity* calculated for a programmer after proportional allocation has been carried out, then it is assumed that they will deliver 25 LoC in the next week (Figure 5e).

The *Achieved Work* is simply the sum of productivity over time. The simple model therefore approximates to linear growth in the product (Figure 5f).

The model is simulated over the period of ten weeks (i.e.  $t=0$  to  $t=10$ ). A manager can then make assessments to see if the loads on each delivery (A1 100 LoC/week and A2 50 LoC/week), and on the code team as a whole ( $A1 + A2 = 150$  LoC/week), are realistic and achievable given the available resources.

		Week									
		1	2	3	4	5	6	7	8	9	10
Inputs	PlannedProduct[D1,Code]	1000									
	PlannedProduct[D2,Code]	500									
	PlannedStartTime[D1,Code]	1									
	PlannedStartTime[D2,Code]	1									
	PlannedEndTime[D1,Code]	10									
	PlannedEndTime[D2,Code]	10									
	PlannedResource[Code]	6									
Outputs	RequiredEffectiveness[D1,Code]	100	100	100	100	100	100	100	100	100	100
	RequiredEffectiveness[D2,Code]	50	50	50	50	50	50	50	50	50	50
	AppliedResource[D1,Code]	4	4	4	4	4	4	4	4	4	4
	AppliedResource[D2,Code]	2	2	2	2	2	2	2	2	2	2
	AppliedEffort[D1,Code]	4	4	4	4	4	4	4	4	4	4
	AppliedEffort[D2,Code]	2	2	2	2	2	2	2	2	2	2
	RequiredProductivity[D1,Code]	25	25	25	25	25	25	25	25	25	25
	RequiredProductivity[D2,Code]	25	25	25	25	25	25	25	25	25	25
	AchievedProduct[D1,Code]	100	200	300	400	500	600	700	800	900	1000
	AchievedProduct[D2,Code]	50	100	150	200	250	300	350	400	450	500



## 5.2 Example 2: Two Staged Deliveries, One Phase

The first example considered two deliveries operating in parallel for the entire duration of the simulation (Example 1). In practice, software deliveries tend to be staged over time in response to the constraints from upstream phases (e.g. the supply of systems requirements), the needs of downstream phases (e.g. the availability of code for engine testing), and to balance the pressure on the development team.

In this example, we consider the impact of staging our two deliveries by delaying the start of *A2* until Week 5, whilst finishing *A1* earlier at Week 8 (Figure 6a). We therefore amend the *Planned Start Time* and *Planned End Time* of the deliveries and run our model again to observe the effects.

The immediate impact of staging the deliveries can be seen in Figure 6b. Having reduced the available timescale for each delivery, whilst keeping the *Available Resource* and *Planned Work* at the same levels, there is a consequent increase in the *Required Effectiveness* of both deliveries.

The effect of the staged deliveries can be seen when we look at the way the *Planned Resource* is allocated across the concurrent deliveries as *Applied Resource* (Figure 6c). All the available resources are applied to delivery *A1* until Week 5 at which time it is split between *A1* and *A2* proportionally relative to the required effectiveness of the deliveries. When delivery *A1* finishes, in Week 8, all the available resource is switched to delivery *A2*. In simple terms, we are therefore modelling the way a manager might allocate their resources and the resulting effort profile (Figure 6d).

The key outputs of the model are the *Required Effectiveness* and *Required Productivity* necessary to perform the given schedule successfully (Figure 6b and Figure 6e). If we compare these results with our earlier example, Figure 5b, we can see that the *Required Effectiveness* has increased from 100 to 125 LoC/Week for *A1* and from 50 to 100 LoC/Week for *A2*. The corresponding demands on *Required Productivity* (Figure 6e) may be relatively trivial and absorbed by a short-term increase in team performance (due to pressure). If the deviation is more significant then there will be a higher risk of schedule

problems and overloading. The manager might therefore choose to (i) adjust the profile of *Planned Resource*, (ii) adjust the *Planned Start Time* and *Planned End Time* of schedule activities, or (iii) adjust the levels of *Planned Work* to be performed in each delivery. The manager can therefore investigate the potential impact of alternative schedules and constraints.

### 5.3 Example 3: Two Staged Deliveries, Two Phases

Until now, we have only been concerned about one lifecycle phase (i.e. *Code*). In practice, we want to model the many levels of concurrency and iteration that we observe in industry. To do this we need to expand our model to consider all phases in the development lifecycle.

In this example, we consider a planned schedule consisting of two deliveries (*A1* and *A2*) each with two lifecycle phases, *Code* and *Low Level Test (LLT)*, over a period of 20 weeks as shown in Figure 7a. This allows us to model the hierarchy of concurrency.

The operation of the model is the same as we have described in our previous examples. The *Required Effectiveness* (Figure 7b) calculates the work rate needed to complete each individual combination of delivery and phase (e.g. *A1 Code*, *A1 Low Level Test*) on time. For simplicity, we assume in the primitive model that resource cannot move between code and test phases, reflecting the skills restrictions observed in industry. The level of *Applied Resource* thus uses independent resource pools for code and test phases, i.e. *Planned Resource (Code)* = 6 and *Planned Resource (LLT)* = 4. At any point, the available coders are allocated among parallel coding phases and the available testers are allocated among parallel testing phases, giving the profile in Figure 7c.

This example illustrates one of the main problems of planning concurrent software development. The manager must ensure that the work-products of upstream phases (e.g. *Code*) are sufficiently mature for downstream phases (e.g. *Low Level Test*) to begin. If the upstream work-products are not sufficiently mature then there may be delays in the schedule. Conversely, if the upstream work-products are “too mature” then we may be missing an opportunity to further compress the schedule. In this example, the code for both products is 50% complete when low level testing begins (as shown by the arrows in Figure 7f). The

manager then has the option to determine if this level is reasonable or not based on their observations on past performance.

## **6. Conclusions**

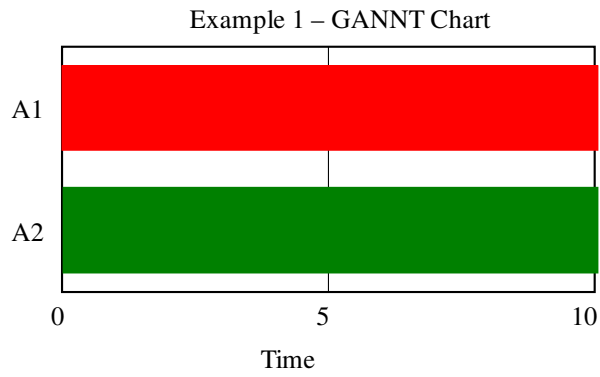
In this paper, we have identified the limitations of current predictive models and proposed a new approach to modelling time-constrained development, called *Capacity-Based Scheduling (CBS)*. The problem with conventional models is that they treat estimation and planning as separate activities. Whereas, in practice, the way a project is planned and controlled has a significant effect on its performance, i.e. '*a different estimate creates a different project*' (Abdel-Hamid 1989).

In response, we have proposed a Capacity-Based Scheduling (CBS) approach to model concurrency and iteration in a planned schedule. A primitive model was used to explain the principles behind the CBS approach. By modelling just two phases of two deliveries, we were able to explain the complexity of interactions in time-constrained projects. A manager must consider simultaneously the effects of planning decisions on concurrent deliveries in the same phase, and subsequent phases of the same delivery, whilst trying to minimise the overall lead-time.

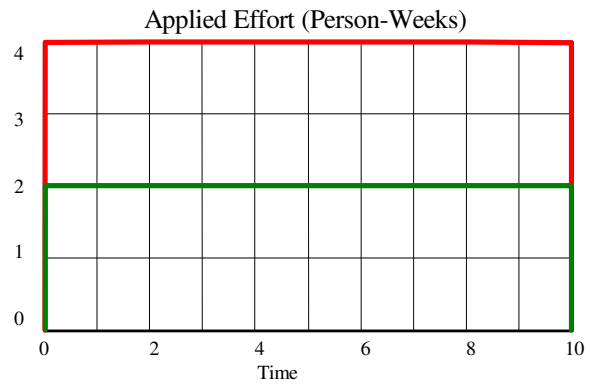
By comparing the levels of required productivity predicted by the model against the actual productivity, it is possible to evaluate the relative feasibility of a plan. The CBS approach therefore allows managers to make much more effective decisions as they can investigate the effect of their choices prior to committing to them.

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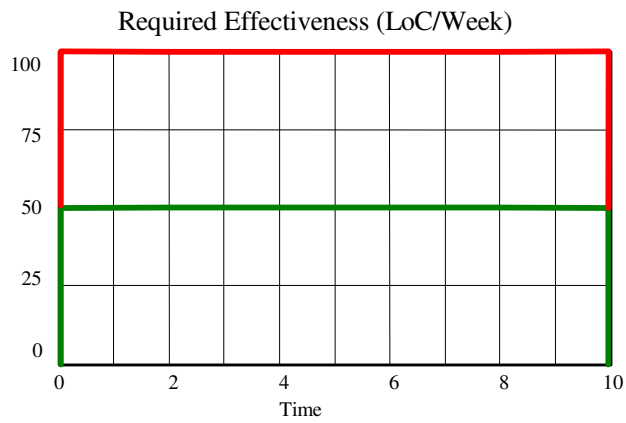
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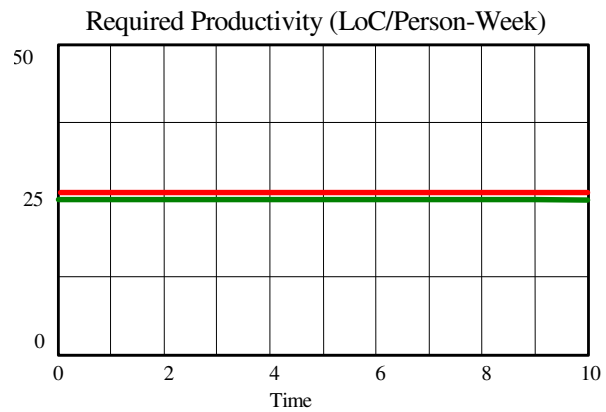
(a)



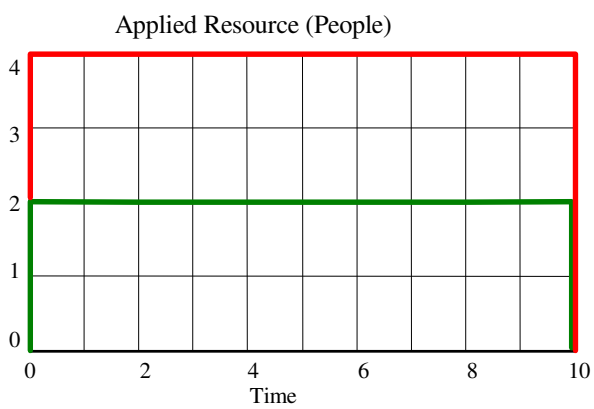
(d)



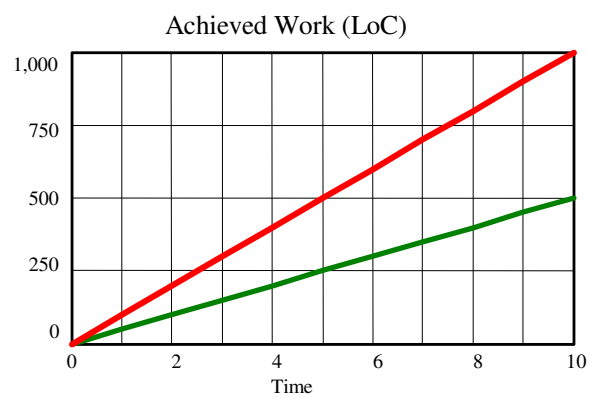
(b)



(e)

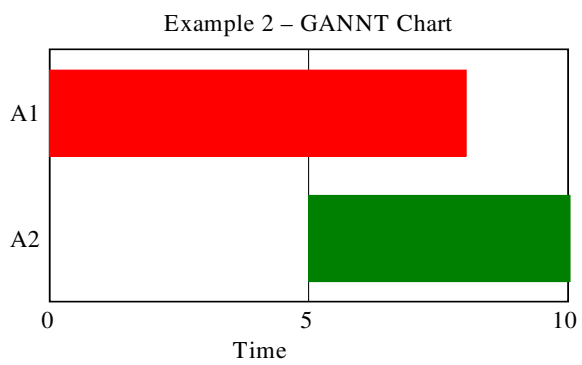


(c)

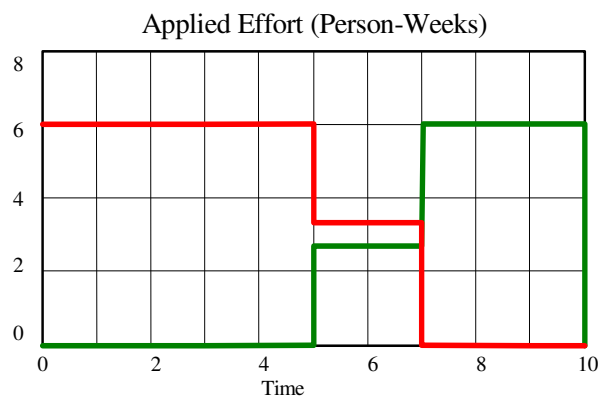


(f)

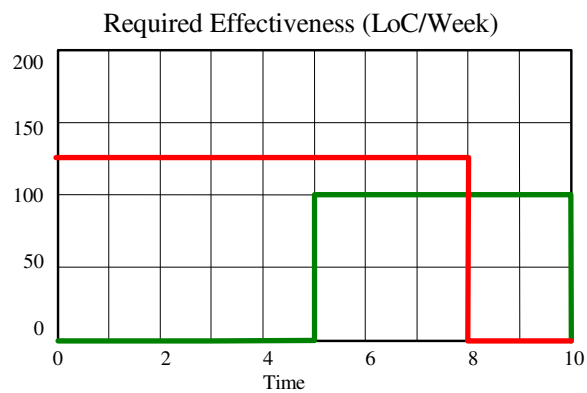
Figure 5: Model Example 1 – Results (One Delivery, One Phase)



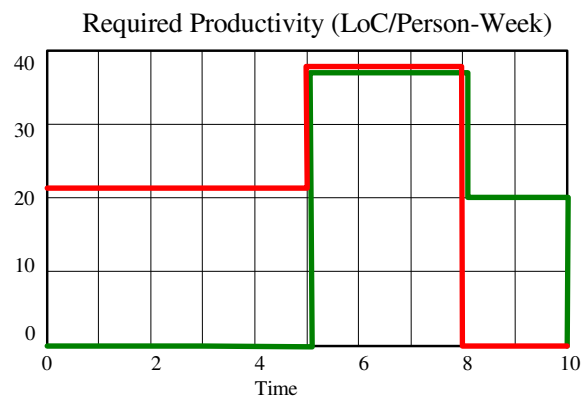
(a)



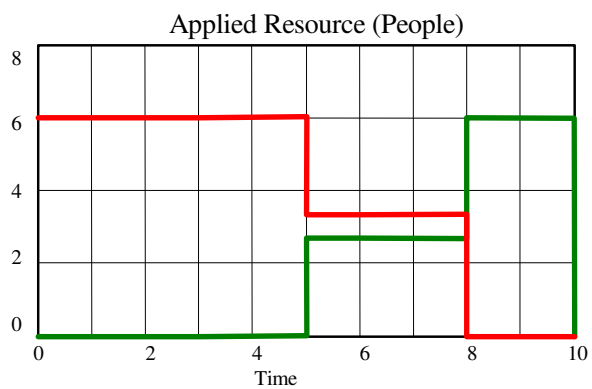
(d)



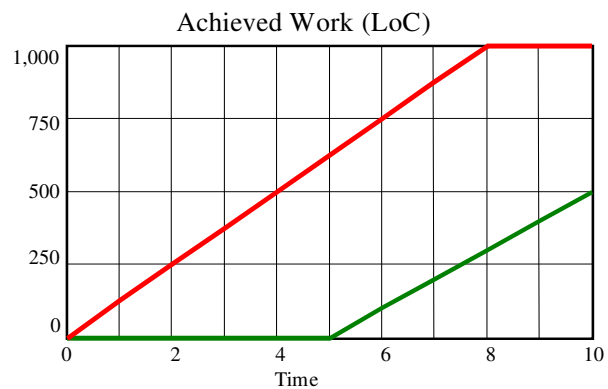
(b)



(e)

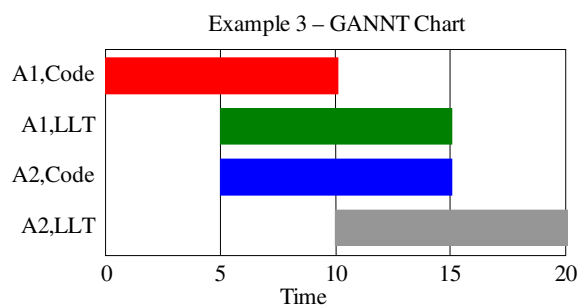


(c)

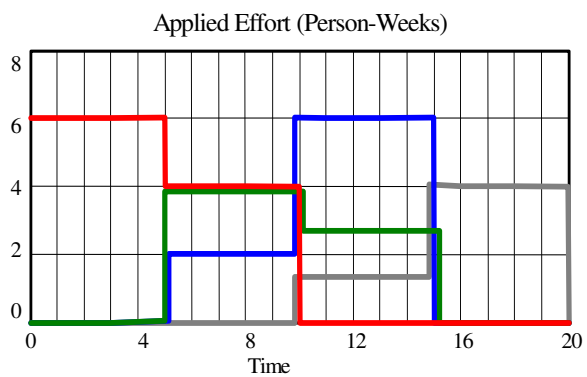


(f)

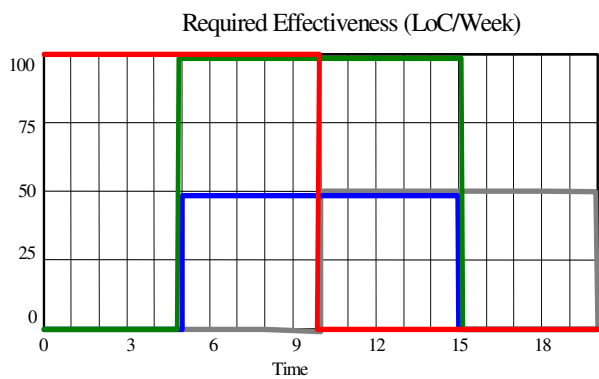
Figure 6: Model Example 2 – Results (Two Deliveries, One Phase)



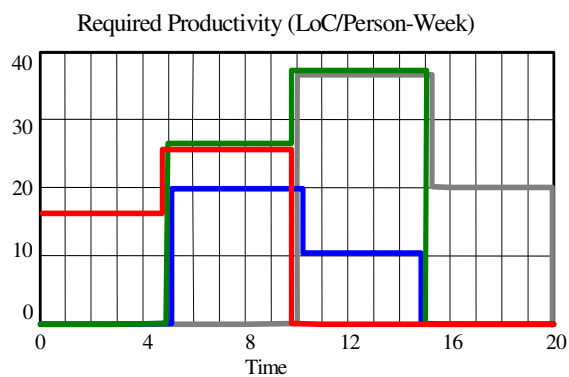
(a)



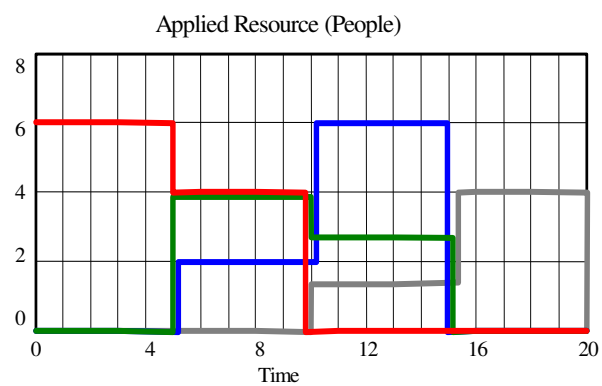
(d)



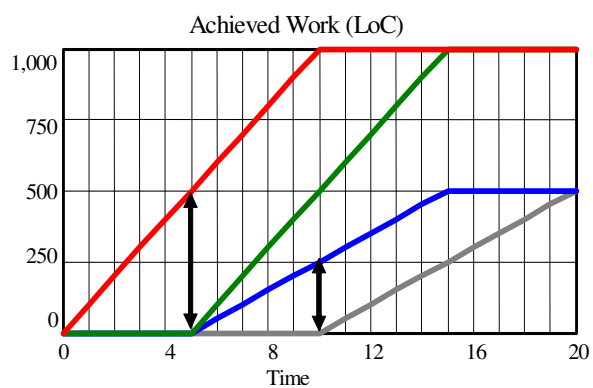
(b)



(e)



(c)



(f)

Figure 7: Model Example 3 – Results (Two Staged Deliveries, Two Phases)